

Review

Thm: Let H be a bipartite on (A, B) such that $v \in A$ has degree at most r . Then $ex(n, H) \leq C n^{2-1/r}$.

Remark: This bound can be tight for $H = K_{r, s}$ where $s > (r-1)!$

Conjecture: Let H be a $K_{r,r}$ -free bipartite graph on (A, B) such that $v \in A$ has degree at most r . Then $\exists c, C > 0$ st. $ex(n, H) \leq C n^{2-1/r-c}$ } open for $r \geq 3$

(*) Today, we will show this conjecture holds for $r=2$.

Thm 1: Let H be a $K_{2,2}$ -free bipartite graph on bipartition (A, B) such that each vertex in A has degree at most 2. Then there exist $c, C > 0$ depending on H such that $ex(n, H) \leq C n^{3/2-c}$

Def: For a graph H , the k -subdivision $H^{(k)}$ of H is a graph obtained from H by replacing each edge ab of H with an internally disjoint path P_{ab} of length $k+1$ with endpoints a and b , where all such paths P_{ab} are mutually internally disjoint.



We point out that all $H^{(k)}$ are bipartite for all graphs H .
(for odd $k \geq 1$)

Thm 2 For all $t \geq 3$, there exists $c_t > 0$ such that $ex(n, K_t^{(t)}) = O(n^{3/2-c_t})$ } $c_t = \frac{1}{4t-6}$

(Caron-Lee, Janzer '19)

We observe that Thm 2 can imply Thm 1. (?)

Because any H in Thm 1 is contained in $K_{|A|}^{(1)}$.
 $\Rightarrow ex(n, H) \leq ex(n, K_{|A|}^{(1)})$

A graph G is called K -regular if $\Delta(G) \leq K \cdot \delta(G)$.

Lemma 1 (Erdős-Simonovits; Jiang, Bukh-Jiang, Caron-Lee)

For all $0 < \alpha < 1$, there exist constants $\beta, K > 0$ such that for all $C > 0$ and sufficiently large n , every n -vertex graph G with at least $C n^{1+\alpha}$ edges has a subgraph G' satisfying

(1) G' is K -regular and bipartite with two parts of sizes differing by a factor at most 2,

(2) $v(G') \geq n^\beta$, and

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(3) $\frac{e(G')}{v(G')^{1+\alpha}} \geq \frac{1}{10} \cdot \frac{e(G)}{v(G)^{1+\alpha}}$

$\Rightarrow e(G') \geq \frac{C}{10} \cdot n^{1+\alpha}$

We will not give a detailed proof for this lemma.

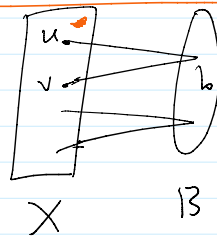
Def. Fix t . For $u, v \in V(G)$, we say the pair $\{u, v\}$ is light, if $1 \leq |N(u) \cap N(v)| < \binom{t}{2}$, and is heavy if $|N(u) \cap N(v)| \geq \binom{t}{2}$.

Lemma 2. Let G be a $K_t^{(1)}$ -free bipartite graph with bipartition $X \cup B$, where $d(x) \geq \delta$ for all $x \in X$ and $|X| \geq \frac{4|B|t}{\delta}$. Then there exists $u \in X$ in $\Omega(\delta^2 |X|/|B|)$ light pairs in X .

Pf. Let $S = \{(\{u, v\}, b) : b \in B, u, v \in X \cap N(b)\}$

We see that

$$|S| = \sum_{b \in B} \binom{d(b)}{2} \geq |B| \binom{e(G)/|B|}{2} \geq \frac{|B|}{4} \left(\frac{\delta |X|}{|B|}\right)^2 = \frac{\delta^2 |X|^2}{4|B|}$$

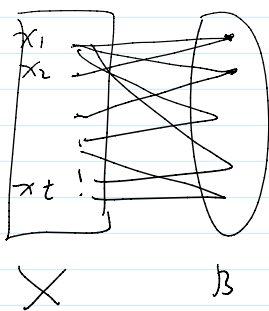


$e(G) \geq \delta |X|$

Let $B^* = \{b \in B : d(b) \geq 2t\}$

Since $\sum_{b \in B \setminus B^*} \binom{d(b)}{2} \leq 2t^2 \cdot |B| \leq \frac{\delta^2 |X|^2}{8|B|}$

$\Rightarrow \sum_{b \in B^*} \binom{d(b)}{2} \geq \frac{\delta^2 |X|^2}{8|B|}$ (1)



Claim: No t vertices in X , any pair of which is heavy

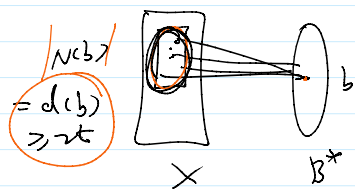
Pf. Suppose say x_1, x_2, \dots, x_t s.t. $\{x_i, x_j\}$ is heavy

i.e. $|N(x_i) \cap N(x_j)| \geq \binom{t}{2}$

It is easy to see that there exists a $K_t^{(1)}$ in G , a contradiction

- Consider $b \in B^*$. Any pair in $N(b)$ is either light or heavy

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By Turán's Thm and the claim, the number of heavy pairs in $N(b)$ is at most $e(T_{t-1}(d(b)))$.

$\Rightarrow \forall b \in B^*$, there are at least

$$\binom{d(b)}{2} - e(T_{t-1}(d(b))) \geq \binom{d(b)}{2} - \binom{t-1}{2} \left(\frac{d(b)}{t-1}\right)^2$$

$$\geq \frac{1}{2(t-1)} d(b)^2 - \frac{1}{2} d(b) \geq \Omega(d(b)^2)$$

light pairs in $N(b)$.

If we sum over all $b \in B^*$, then

$$\# \sum_{\substack{\{u,v\} \text{ is light} \\ \text{in } N(b) \text{ \& } b \in B^*}} \binom{\{u,v\}, b} \geq \sum_{b \in B^*} \Omega(d(b)^2)$$

$$\stackrel{(1)}{\geq} \Omega\left(\frac{\sum |X|^2}{|B|}\right)$$

Since $\{u,v\}$ is light, we get

$$\# \text{ light pairs in } X \geq \frac{\#}{\binom{t}{2}} \geq \Omega\left(\frac{\sum |X|^2}{|B|}\right)$$

$\Rightarrow \exists$ a vertex $u \in X$ which is in at least

$$\Omega\left(\frac{\sum |X|^2}{|B|}\right) \text{ light pairs in } X. \quad \square$$

Pf of Thm 2. We give a proof due to Janson,

who proved that $c_t = \frac{1}{4t-6}$ for $t \geq 3$.

$t=3$ $K_3^{(1)}$ = C_6

perhaps $c_t = \frac{1}{4t-6}$

$ex(n, C_6) = \Theta(n^{4/3})$ is tight; is tight for all $t \geq 3$ (?)

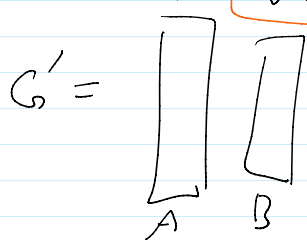
Let G be a $K_t^{(1)}$ -free graph on n vertices & with

at least $D \cdot n^{\frac{3}{2}-\epsilon} = D n^{1+\alpha}$ ($\alpha = \frac{t-2}{2t-3}$) edges.

By Lemma (with α), there exists a $G' \subseteq G$

l.i.l is $\lfloor \frac{1}{2} \log \frac{1}{\epsilon} \rfloor$...

by Lemma 1 (with α), there exists a $G' \subseteq G$
 which is K -regular and bipartite on parts $A \cup B$,
 such that $\frac{e(G')}{v(G')^{1+\alpha}} \geq \frac{1}{\omega} \frac{e(G)}{v(G)^{1+\alpha}}$ & $v(G')$ is large,



$$\frac{1}{2}|B| \leq |A| \leq 2|B|$$

⊗ If $\delta(G') \leq C \cdot (v(G'))^\alpha$
 then $\Delta(G') \leq K \cdot C \cdot (v(G'))^\alpha$
 $\Rightarrow e(G') \leq K C \cdot (v(G'))^{1+\alpha}$

$$\Rightarrow e(G) \leq \omega K C \cdot n^{1+\alpha} \quad \text{done}$$

Therefore, $\sum \delta(G') \geq C \cdot (v(G'))^\alpha = C \cdot (v(G))^{(t-2)/(t-3)}$

Our plan is to find t vertices u_1, u_2, \dots, u_t in A
 such that $\{u_i, u_j\}$ is light $\forall 1 \leq i < j \leq t$.

* u_i, u_j, u_k have no common neighbors for all distinct i, j, k .

If so, then we can find a $K_t^{(1)}$ in G easily.

We will do so by repeatedly using Lemma 2 on a stronger

hypothesis: for each $1 \leq i \leq t$, there exists $A = X_1 \supseteq X_2$
 $\supseteq X_3 \dots \supseteq X_i$ and $u_1 \in X_1, u_2 \in X_2, \dots, u_i \in X_i$ s.t.

(a) u_j is in at least $\Omega(\delta(X_j)/v(G))$ light pairs in X_j
 for $\forall 1 \leq j \leq i-1$.

(b) u_j is light to every vertex w in $X_{j+1} \Leftrightarrow \{u_j, w\}$ is light
 $\forall 1 \leq j \leq i-1$

(c) no 3 of v_1, \dots, v_i have common neighbors

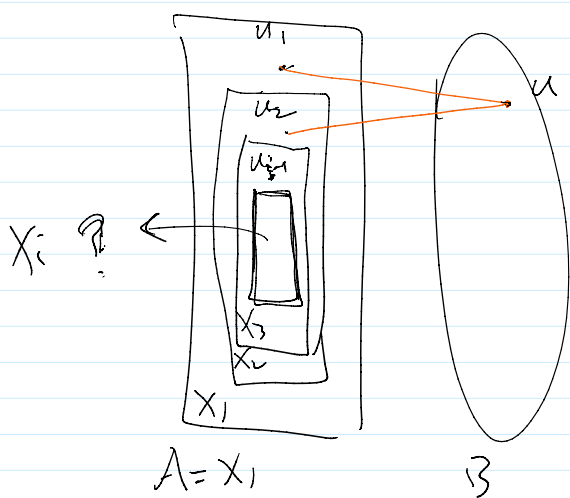
(d) $|X_{j+1}| = \Omega(\delta(X_j)/v(G)) \quad \forall 1 \leq j \leq i-1$

This holds clearly for $i=1$ by choosing u_1 to be the vertex
 found by Lemma 2 when applied to $A \cup B$.

Now suppose we have constructed this for $i-1$.

$A = V \supseteq X_1 \supseteq \dots \supseteq X_{i-1}$ with $u_i \in X_{i-1}$

Now suppose we have constructed this for $i-1$:
 $A = X_1 \supseteq X_2 \supseteq \dots \supseteq X_{i-1}$ with $u_j \in X_j \forall j \leq i-1$.



Let $Y_i = \left\{ y \in X_{i-1} : \{y, u_{i-1}\} \text{ is light} \right\}$

By (a), $|Y_i| \geq \delta^2 |X_{i-1}| / v(G')$

Consider any u_j, u_ℓ , $j, \ell \leq i-1$, take any common neighbor u of them and delete $N(u)$ from Y_i

There are $\binom{i-1}{2}$ pairs u_j, u_ℓ & there are at most $\binom{t}{2}$ many choices of u for each u_j, u_ℓ & there are at most $\leq \delta$ vtx in $N(u)$.

\Rightarrow # vtx deleted in this round is at most $\leq \binom{i-1}{2} \binom{t}{2} \leq \delta = O(\delta)$.

As long as $|Y_i| \geq \delta^2 |X_{i-1}| / v(G') \geq \Omega(\delta)$, then we can get a $X_i \subseteq Y_i$ of size at least

$|X_i| \geq \Omega\left(\frac{\delta^2}{v(G')}\right) |X_{i-1}|$, which satisfies (c).

This is true, because $i \leq t$.

$$\left(\frac{\delta^2}{v(G')}\right)^{i-1} |A| \geq \left(\frac{\delta^2}{n}\right)^{i-1} \cdot n \geq \Omega(\delta)$$

$$\Leftrightarrow \delta^{2t-3} \geq n^{t-2}$$

This shows that the algorithm can keep going until we have $X_1 \supseteq X_2 \supseteq \dots \supseteq X_t$

\forall n there is $\delta = \delta(n) > 0$ such that $1 < i < t$.

going until we have $n_1 - n_2$
and vertices $u_j \in X_j, 1 \leq j \leq t$.

This is clear from this to see

• $\{u_i, u_j\}$ is light, $i \neq j$

• any 3 of u_1, u_2, \dots, u_t have no common neighbors

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